A Walk on the Wild Side: a Shape-First Methodology for Orthogonal Drawings



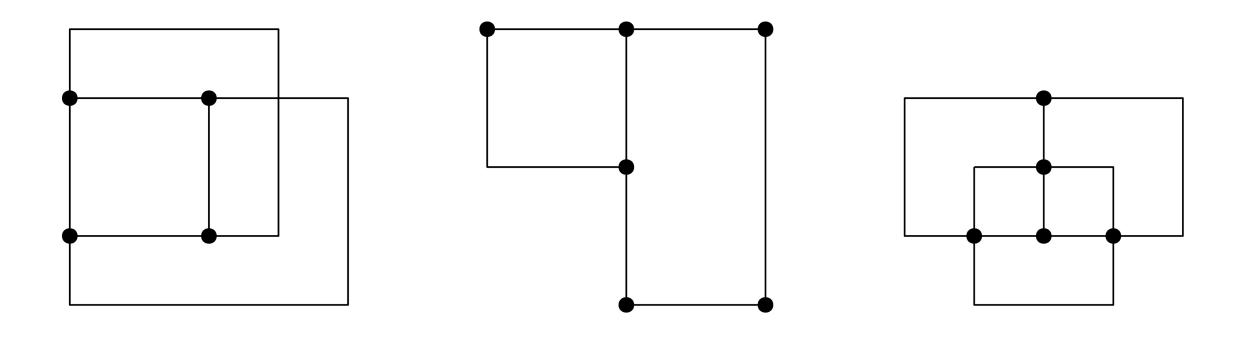
speakers: G. Andreola and A. Strippoli

joint work with: S. Caroppo, G. Di Battista,

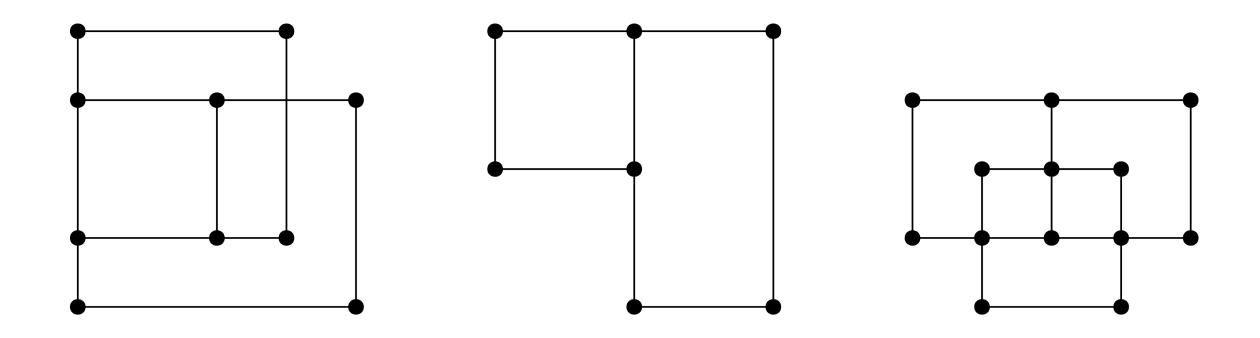
F. Grosso, and M. Patrignani



orthogonal drawing



rectilinear drawing



orthogonal drawings tools





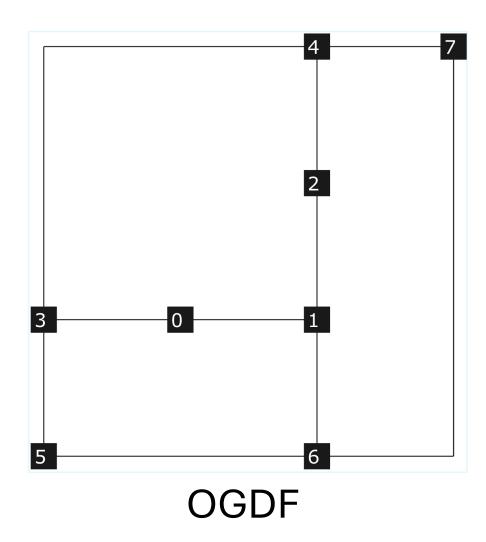




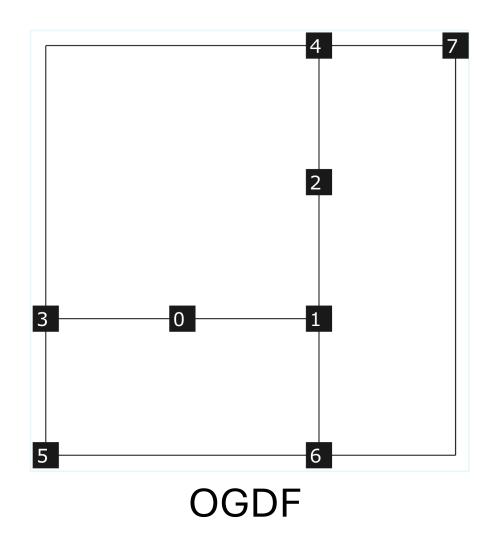


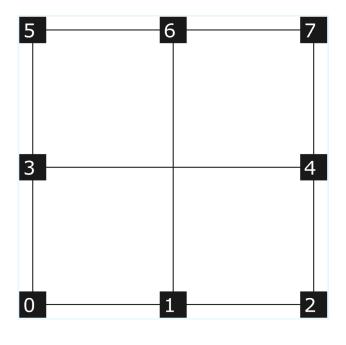


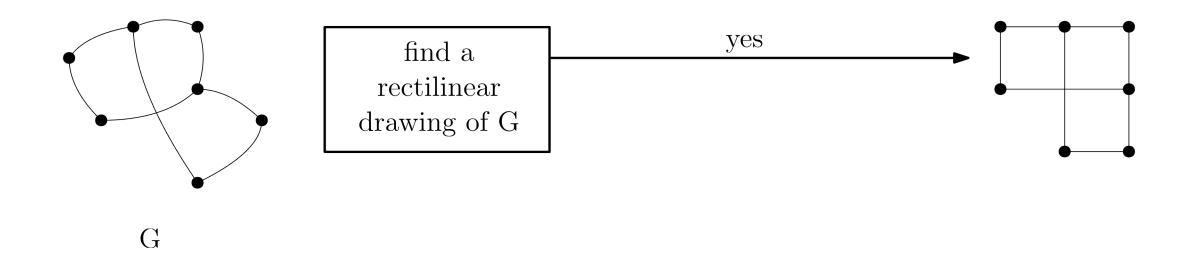
they prioritize crossings...

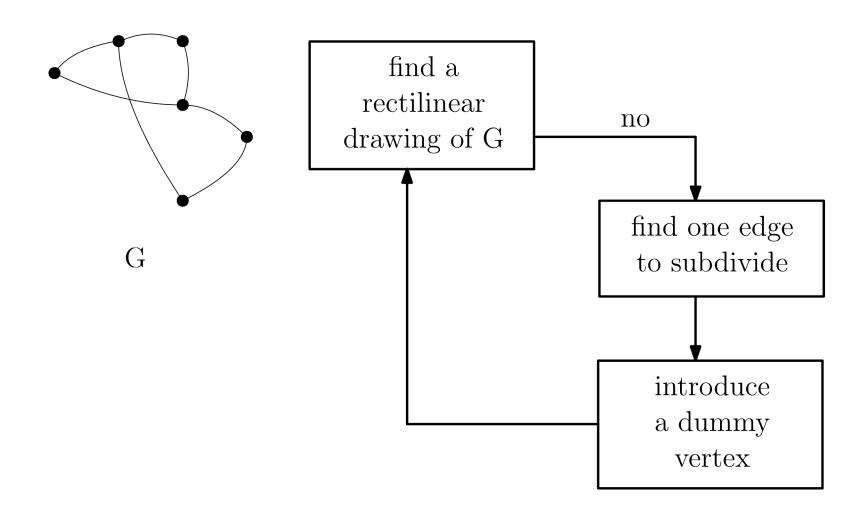


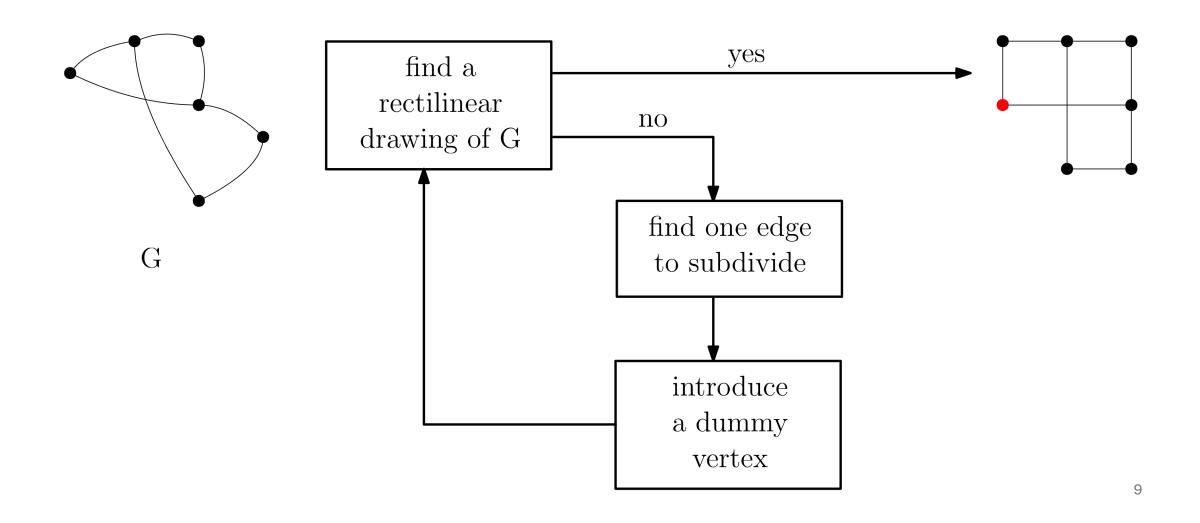
...we prioritize bends

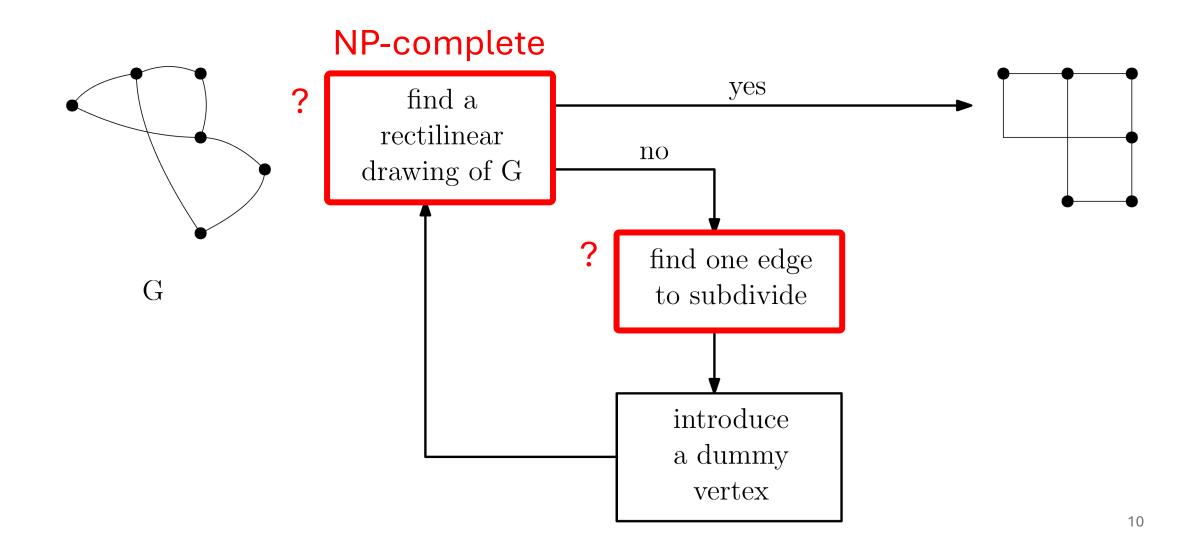








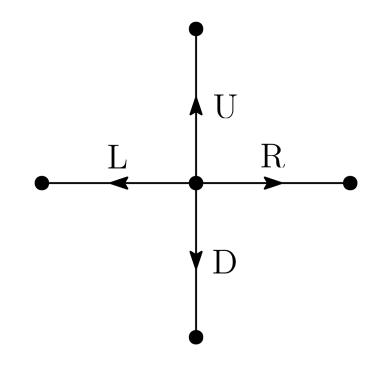




polynomial algorithm for shaped graphs

theorem: given a graph, if each edge has prescribed direction "left", "right", "down" or "up", there exists a polynomial time algorithm that finds a rectilinear drawing of it, if it exists

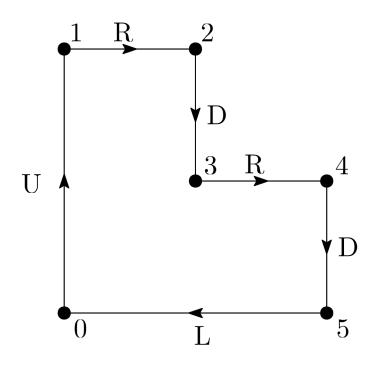
[Manuch, Patterson, Poon, and Thachuk, 2010]

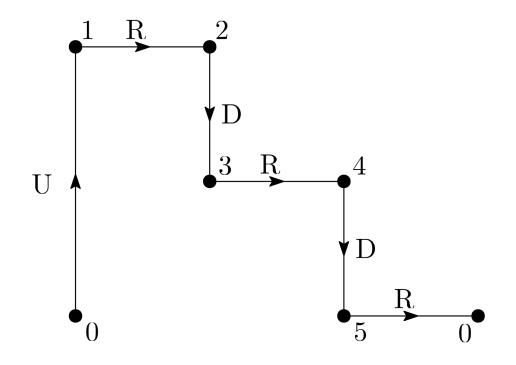


shaped graph

completeness of cycles

definition: a cycle is complete if it contains all the four directions



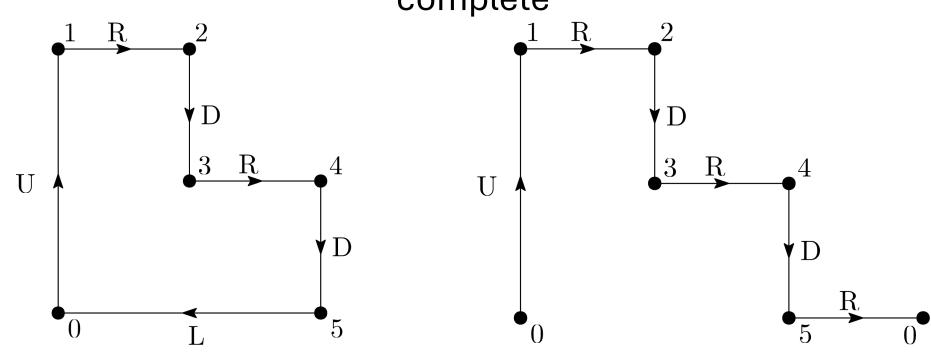


complete cycle

non complete cycle

characterization of rectilinear drawability

theorem: a shaped simple cycle is rectilinear drawable if and only if it is complete

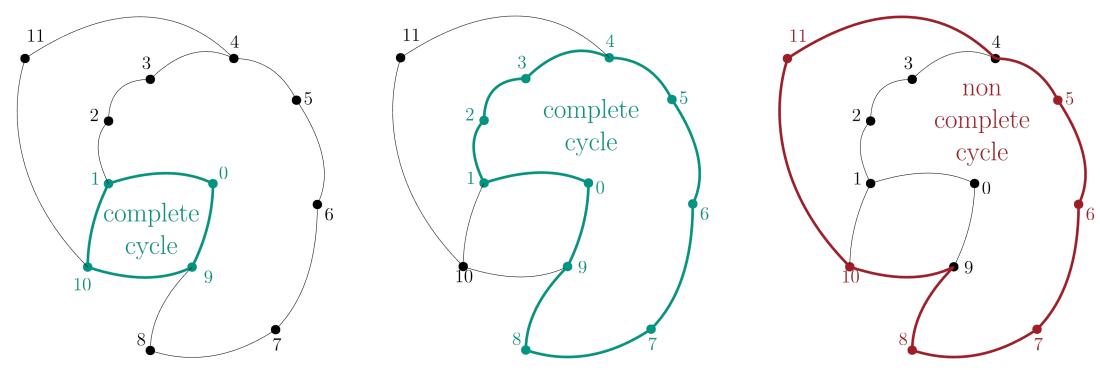


complete cycle

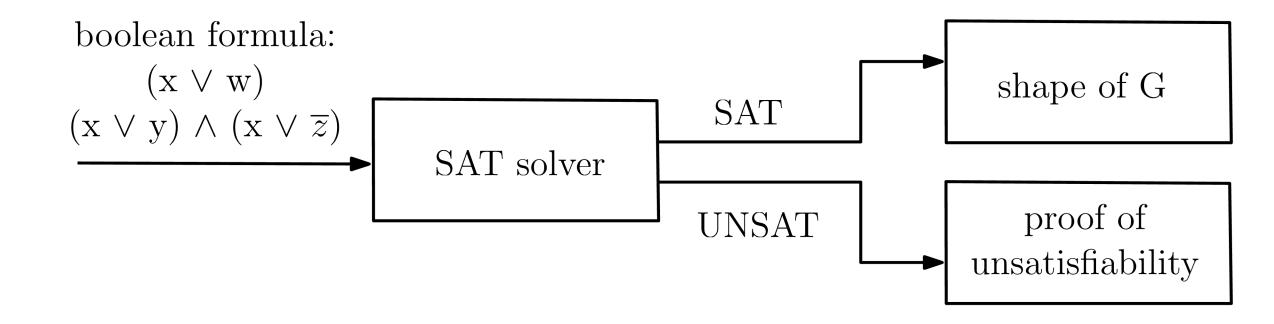
non complete cycle

characterization of rectilinear drawability

theorem: a shaped graph is rectilinear drawable if and only if all its simple cycles are complete



SAT solver for finding a shape



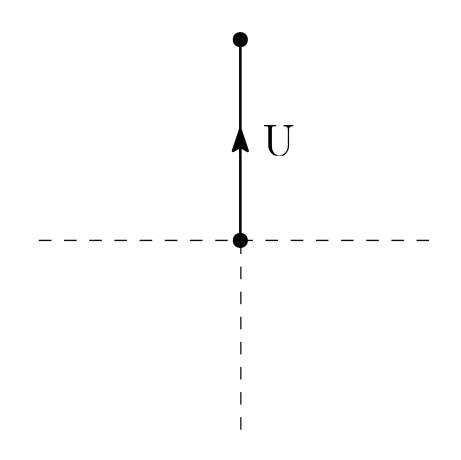
variable	value
R	False
D	False
U	False
\mathbf{L}	False

variable	value
R	True
D	False
U	False
L	False

variable	value
R	False
D	True
U	False
L	False

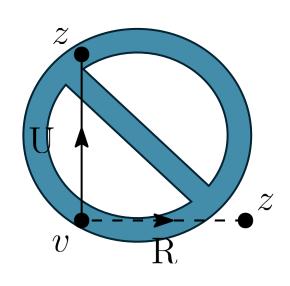
variable	value
R	False
D	False
U	False
L	True

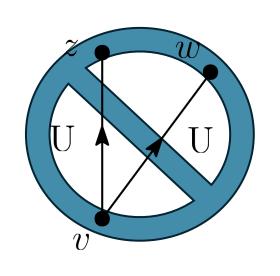
variable	value
R	False
D	False
U	True
L	False

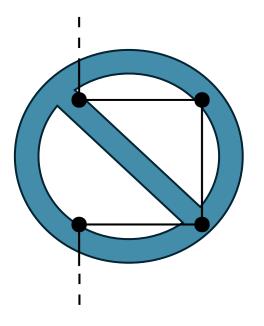


$$\#$$
variables = $4|E| = O(|V|)$

SAT formulation: clauses



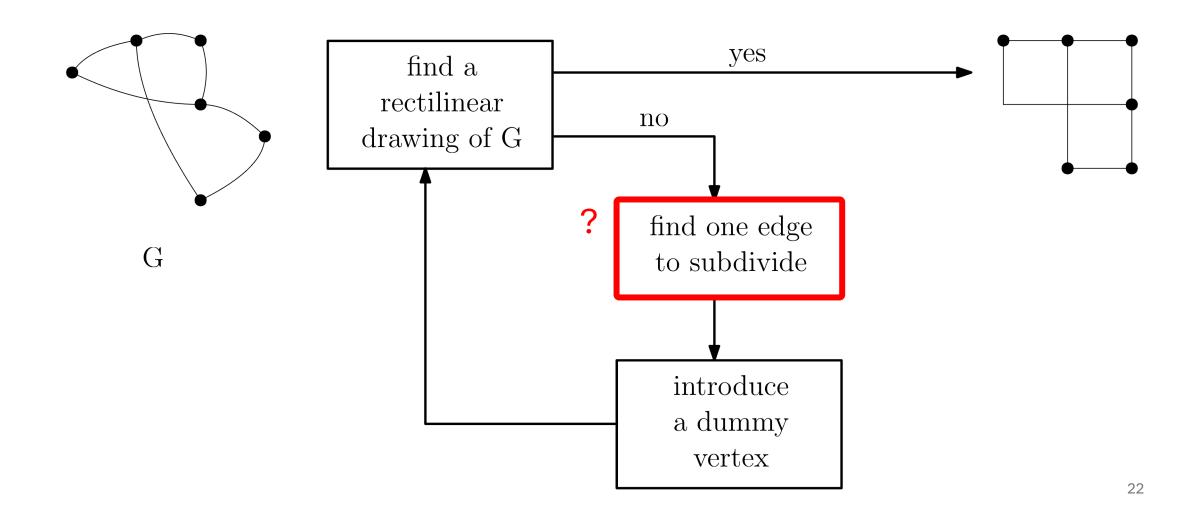




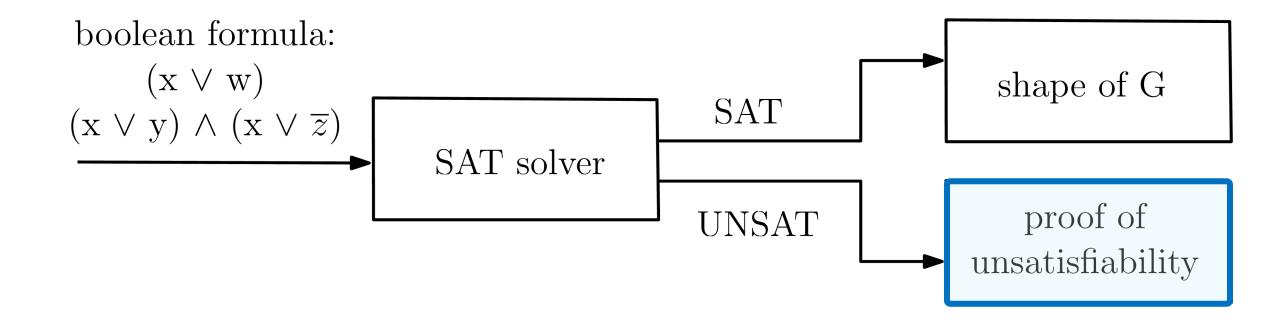
$$\#$$
clauses = $O(|V|)$
clauses length = $O(1)$

$$\#$$
clauses $= O(|C|)$

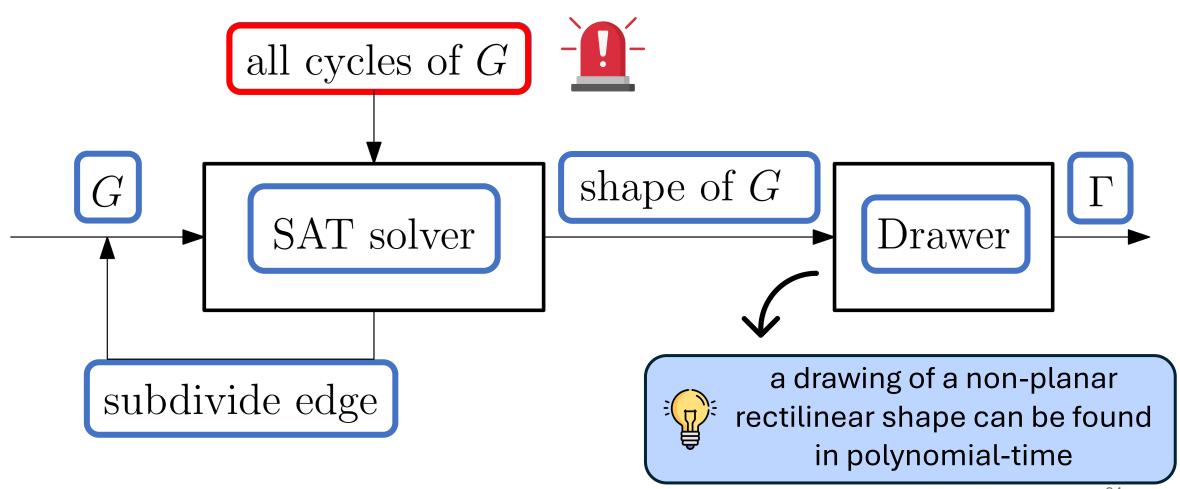
clauses length = O(longest cycle length)



SAT solver for finding a shape

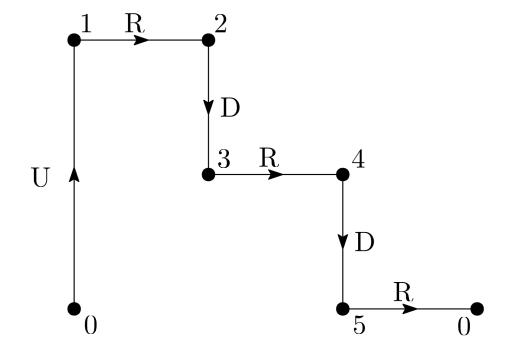


a first framework

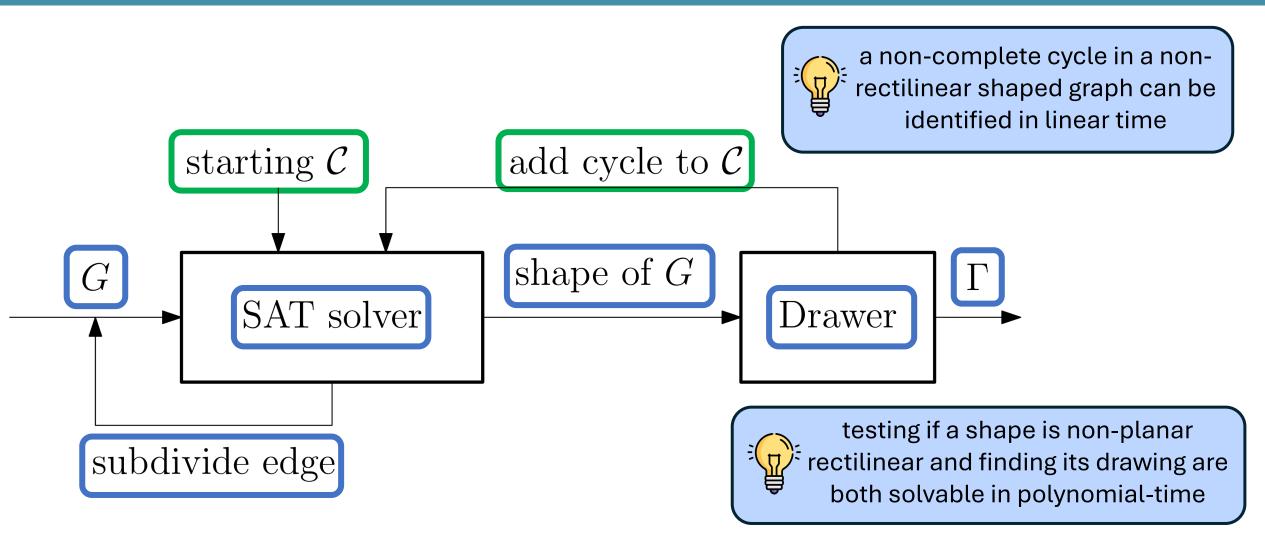


a better approach

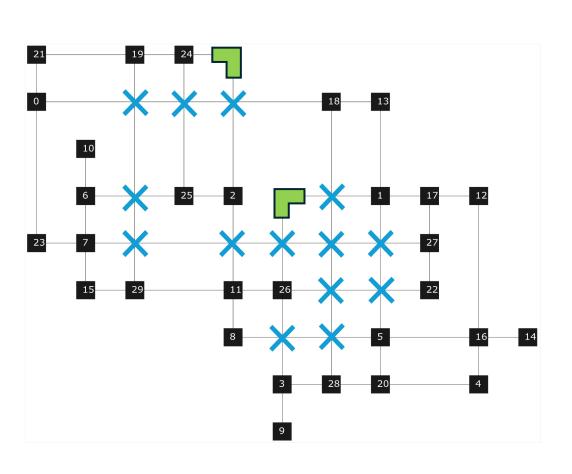
theorem: given a shaped graph that is not rectilinear drawable, a non-complete cycle can be identified in linear time

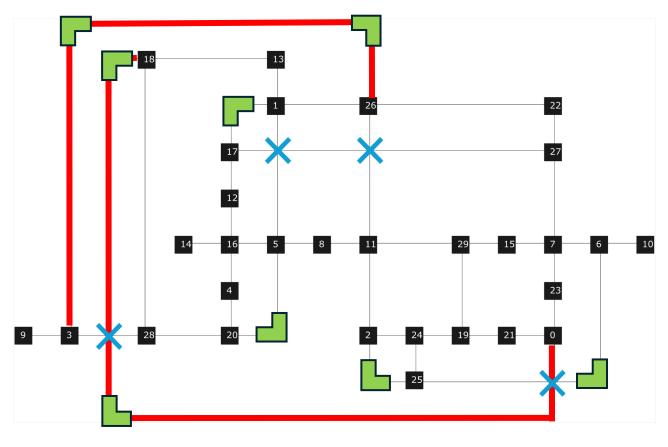


the complete framework



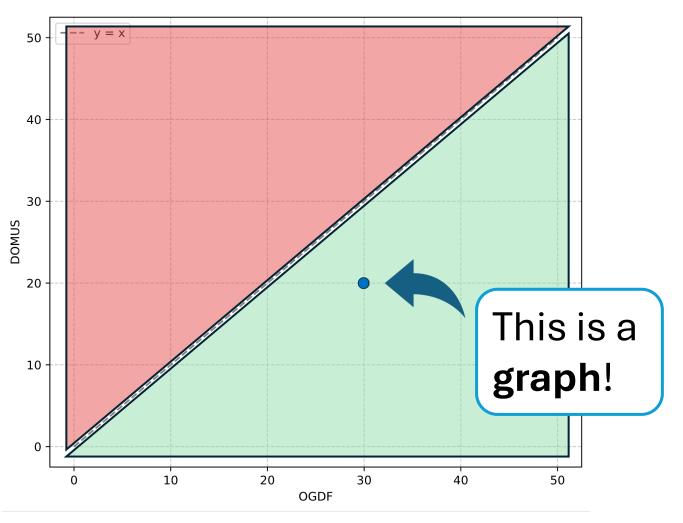
comparing frameworks





DOMUS

OGDF

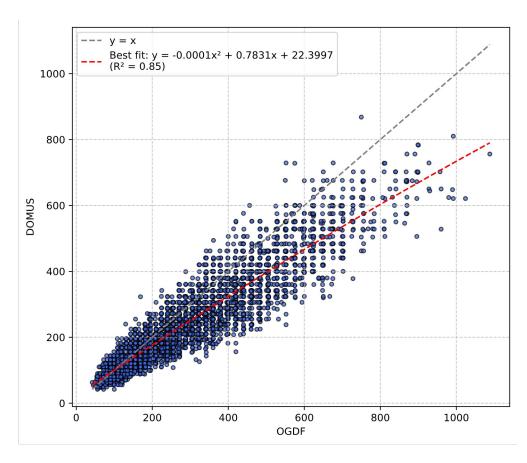


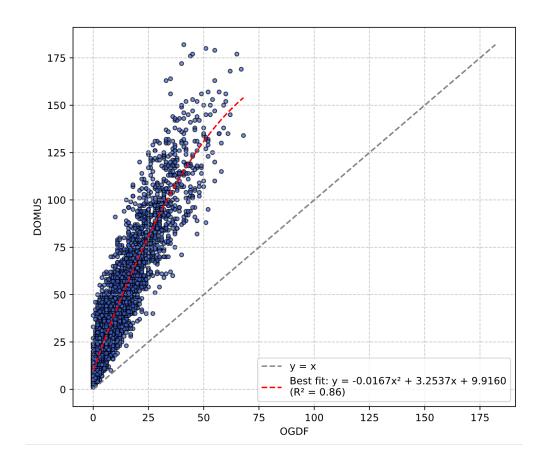
• 4100 graphs

• |V| from 20 to 60

• |E|/|V| from 1.25 to 1.75

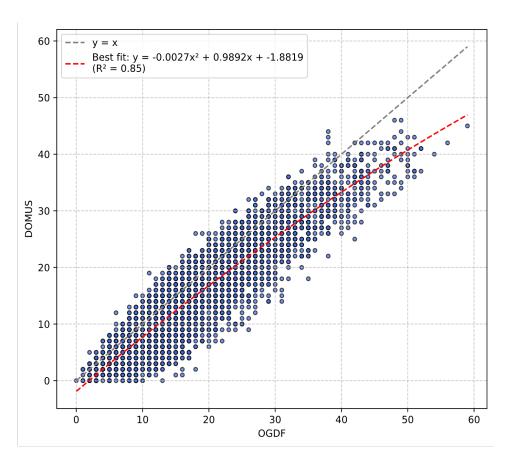






total area

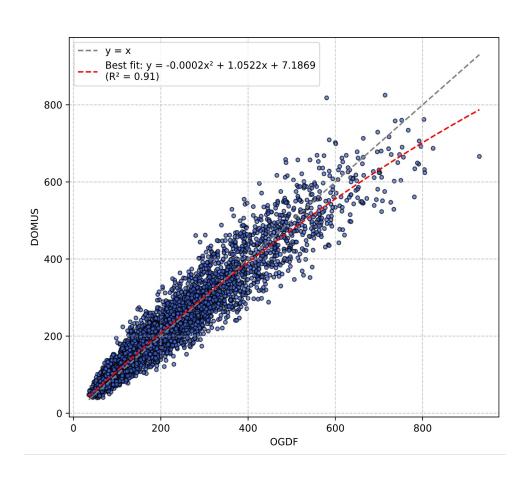
crossings



Best fit: $y = -0.7307x^2 + 1.6178x + -0.2095$ (R² = 0.54) 0.8 0.6 0.2 0.0 0.2 0.8 **OGDF**

total bends

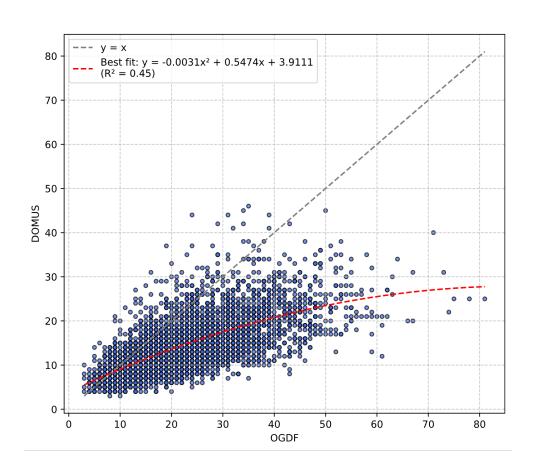
bends std dev

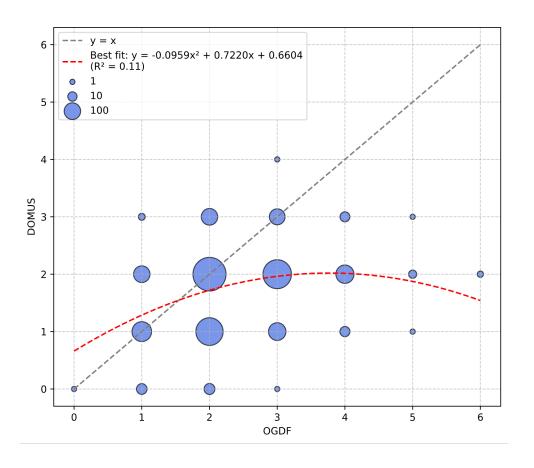


Best fit: $y = -0.0100x^2 + 0.5913x + 0.7530$ 10 DOMUS 10 12 **OGDF**

total edge length

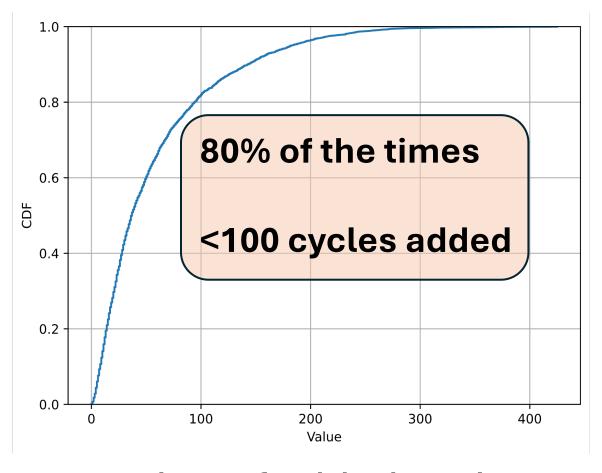
edge length std dev



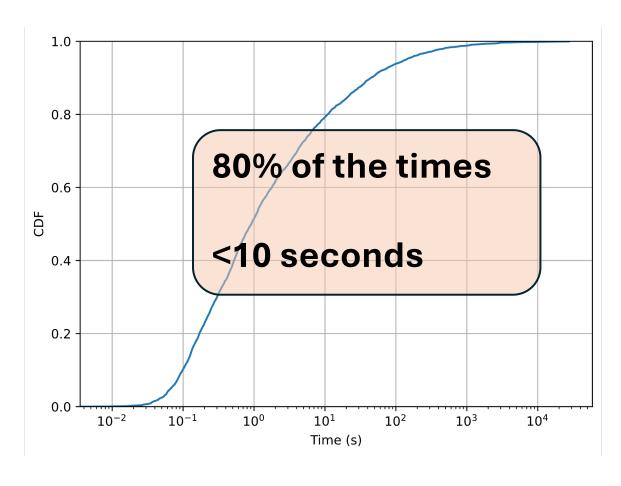


max edge length

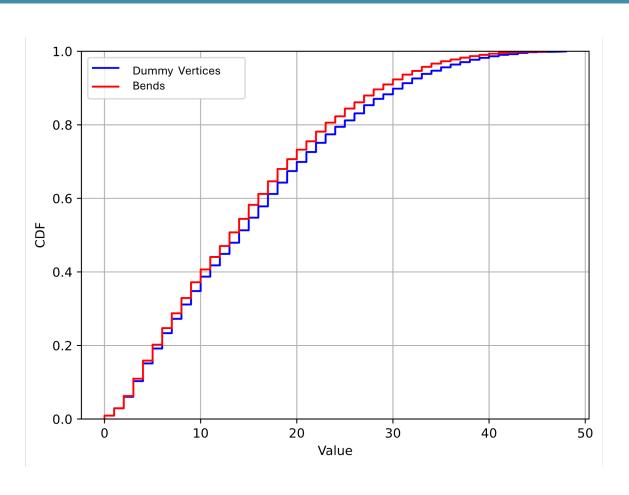
max bends per edge

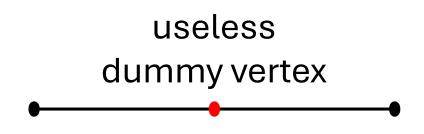


number of added cycles



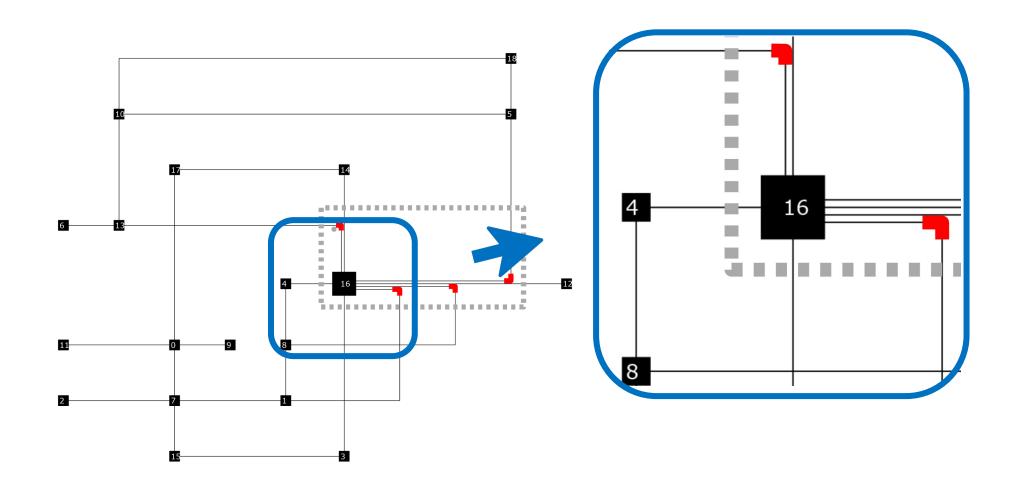
total computation time (seconds)





dummy vertices vs bends

generalizing to any graph



conclusions and open problems

conclusions

- new methodology for orthogonal/rectilinear drawing
- methodology SIMPLE
- good results in most metrics
- constraints can easily be encoded

open problems

- how many cycles to guarantee rectilinear drawability
- encoding crossings in the formulation

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